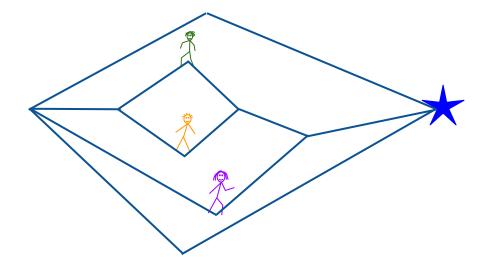
Informational Substitutes



Yiling Chen Bo Waggoner

Harvard UPenn TCS+ talk Nov 9, 2016

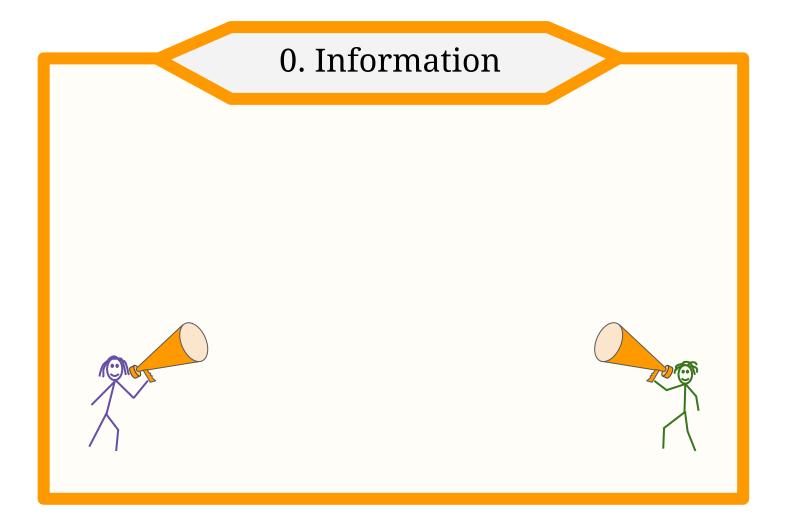
Goals for this talk

1. Introduce a theory audience to some fun topics:

- value of information
- proper scoring rules
- prediction markets

2. Describe this paper (Chen, Waggoner FOCS 2016):

- propose definitions of subs. and comps.
- characterize "good" and "bad" equilibria of prediction markets
- connect to complexity of information acquisition

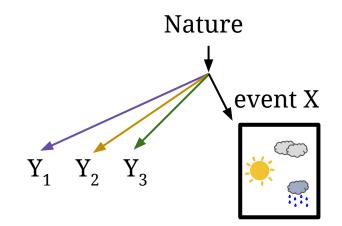


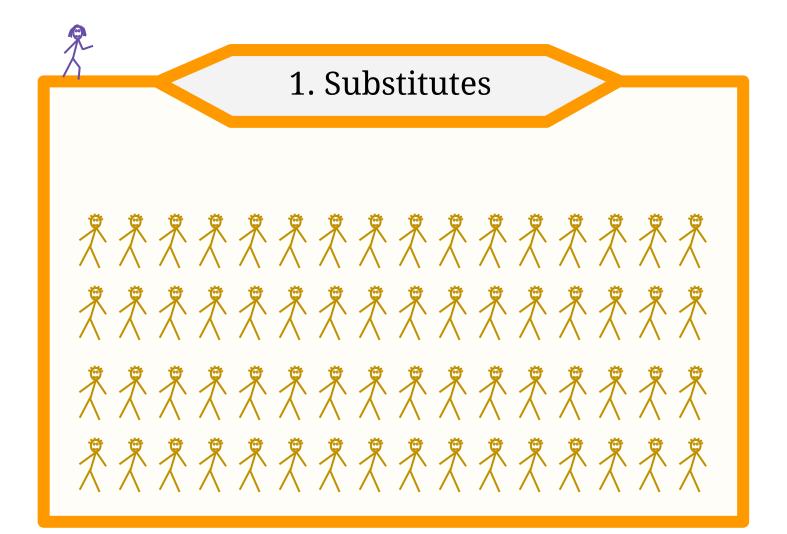
Information (in this talk)

Random variables X, Y_1 , ..., Y_n jointly distributed, known prior. (finite set of outcomes)

We care about X.

Y_i = "signal" (reveals info. about X).





The unreasonable effectiveness of substitutes

Substitutes in economics:

- Market equilibria, stable matchings, ...
- [Kelso & Crawford 1982, Roth 1984, Hatfield and Milgrom 2005, ...]

Substitutes in computer science:

- Submodularity! [Lehmann+Lehmann+Nisan 2001]
- subs == efficient approx. for **many** problems

Could we also define "substitutes" for information?

And could they also link algorithms and game theory?

Challenges for defining informational S&C

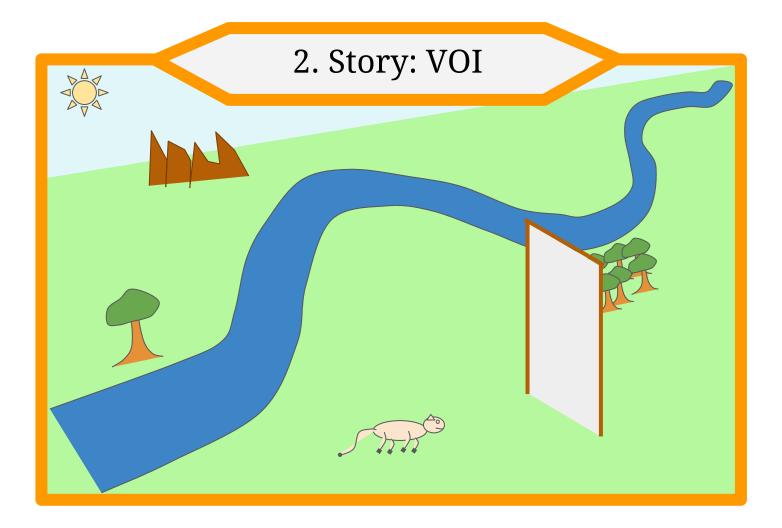
<u>Items</u>

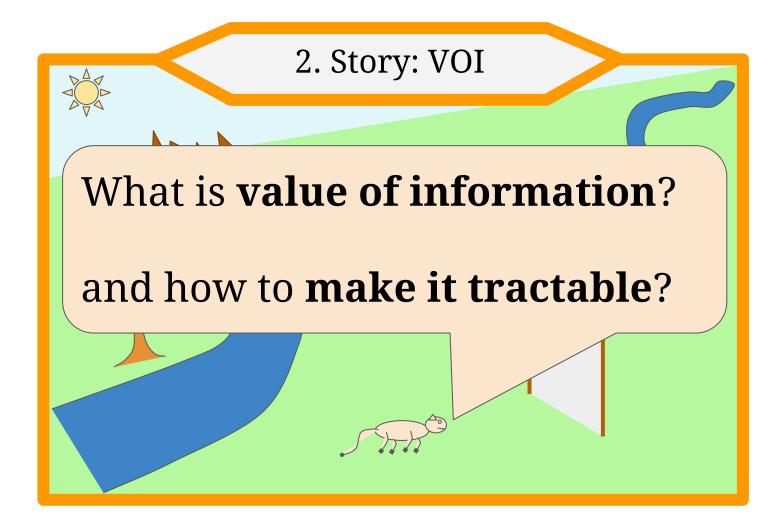
• Valuation function f is given

- f(i) does not depend on f(j)
- "marginal value" is straightforward

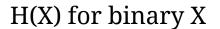
Information

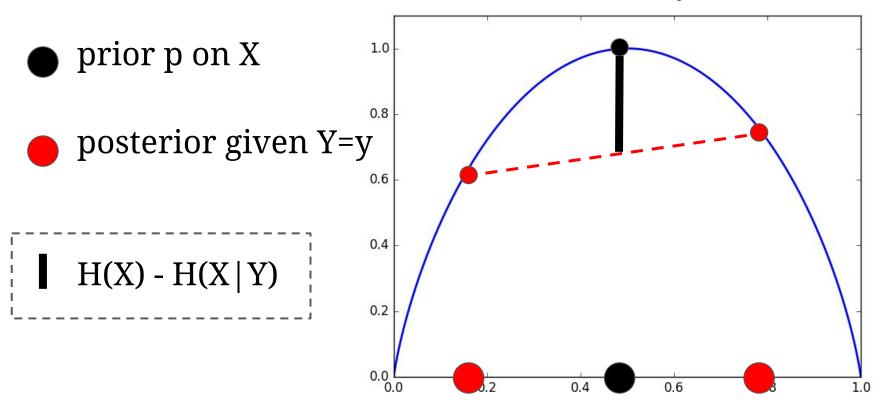
- What is the "value" of a set of pieces of information?
- Two pieces of information may be correlated, redundant, ...
- What is a "marginal" piece of information?





The Story part 1: Shannon 1948





Example where Shannon does not apply

X = rain or no rain Y = weather forecast

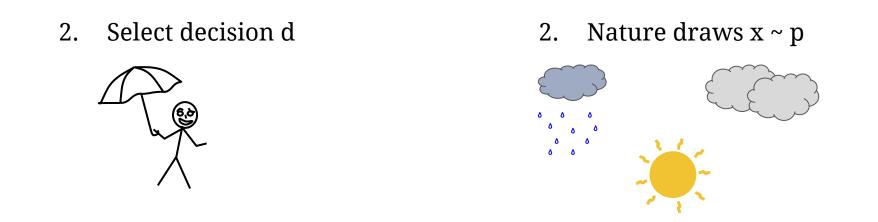
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I hate carrying my umbrella.... But if it rains on me \rightarrow I melt

The Story part 2: Howard 1966

1. Known prior p on X



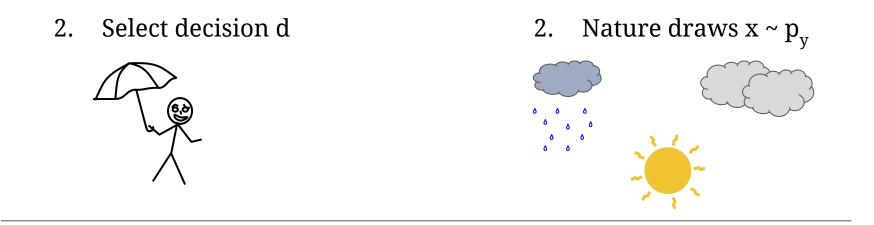
3. Get utility u(d, x).

V(Ø) = "expected utility when deciding optimally with **no signals**"

The Story part 2: Howard 1966

1. Known prior p on X

1.5. **Observe Y**, Bayesian update to p_v



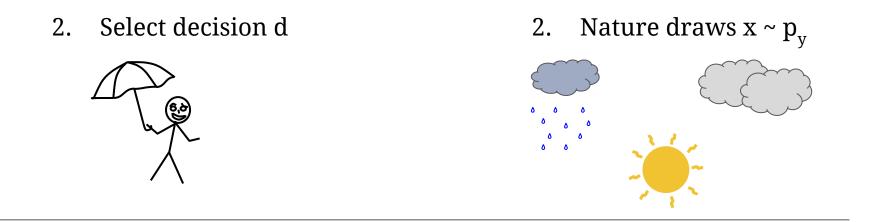
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V(Y) = "expected utility when deciding optimally after **observing** Y"

The Story part 2: Howard 1966

1. Known prior p on X

1.5. **Observe Y**, Bayesian update to p_v



3. Get utility u(d, x).

V(Y) = "expected utility when deciding optimally after **observing** Y"

V(Y) - V(∅) = "marginal value of Y"

But: how to make VOI tractable?

Our approach: apply ideas from proper scoring rules

DEPARTMENT OF COMMERCE CHARLES SAWYER, Secretary WEATHER BUREAU F. W. REICHELDERFER, Chief

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VERIFICATION OF FORECASTS EXPRESSED IN TERMS OF PROBABILITY

GLENN W. BRIER

U. S. Weather Bureau, Washington, D. C. [Manuscript received February 10, 1950]

INTRODUCTION

Verification of weather forecasts has been a controversial subject for more than a half century. There are a number of reasons why this problem has been so perplexing to meteorologists and others but one of the most important difficulties seems to be in reaching an agreement on the specification of a scale of goodness for weather forecasts. Numerous systems have been proposed but one of the greatest arguments raised against forecast verification is that forecasts which may be the "best" according to the accepted system of arbitrary scores may not be the most useful forecasts. In attempting to resolve this difficulty the forecaster may often find himself in the position of choosing to ignore the verification system or to let it do the forecasting for him by "hedging" or "playing the system." This may lead the forecaster to forecast something other than what he thinks will occur, for it is often easier to analyze the effect of different possible forecasts on the verification score than it is to analyze the weather situation. It is generally agreed that this state of affairs

numerically have been discussed previously [1, 2, 3, 4] so that the purpose here will not be to emphasize the enhanced usefulness of such forecasts but rather to point out how some aspects of the verification problem are simplified or solved.

VERIFICATION FORMULA

Suppose that on each of n occasions an event can occur in only one of r possible classes or categories and on one such occasion, i, the forecast probabilities are f_{a} , f_{a} , \ldots , f_{tr} , that the event will occur in classes 1, 2, \ldots , r, respectively. The r classes are chosen to be mutually exclusive and exhaustive so that

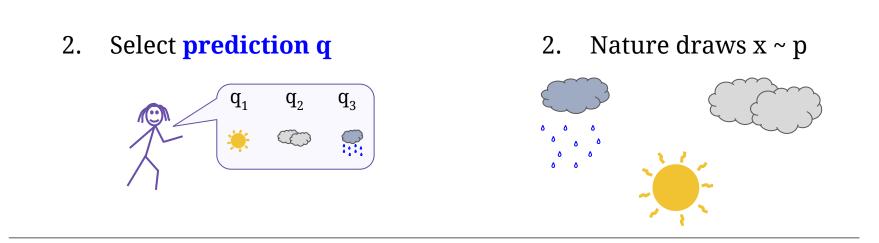
$$\sum_{i=1}^{r} f_{ij} = 1, \, i = 1, 2, 3, \dots n \tag{1}$$

A number of interesting observations can be made about a vertification score P defined by

$$P = \frac{1}{2} \sum_{n=1}^{r} \sum_{n=1}^{n} (f_{n} - F_{n})^{2}$$
(2)

The Story part 3: Savage 1971, "scoring rules"

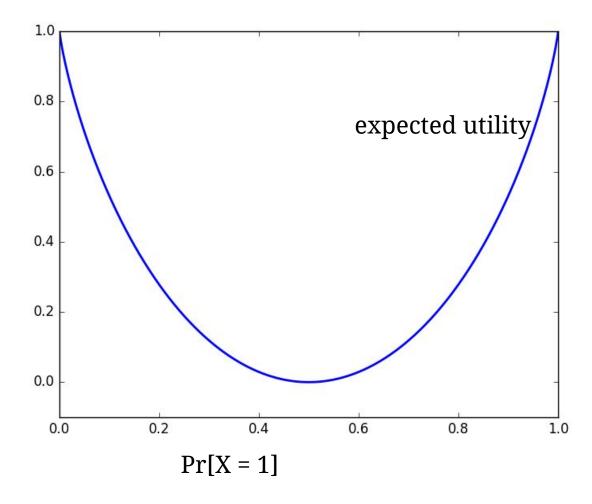
1. Known prior p on X



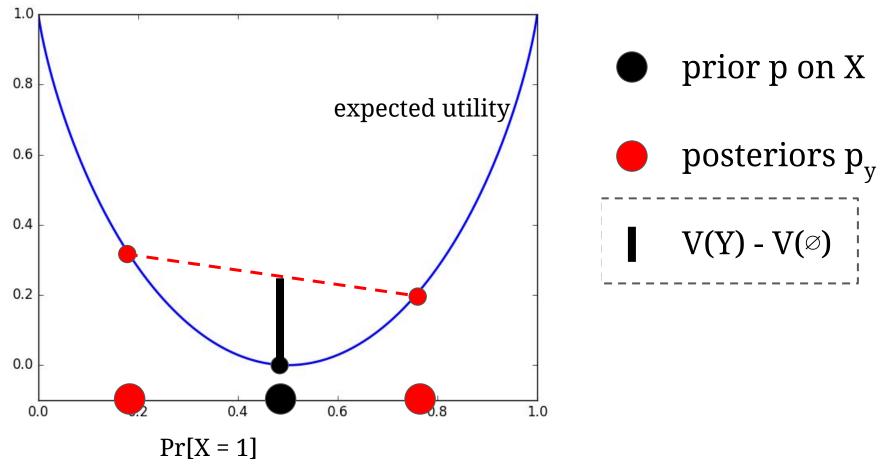
3. Get utility **S(q, x)**.

"Proper scoring rule" - optimal prediction is true belief

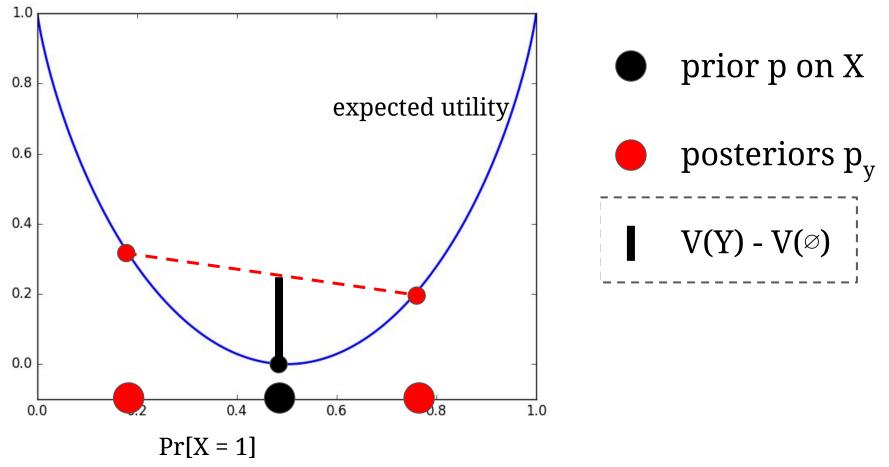
Savage: scoring rules $\leftarrow \rightarrow$ convex functions!

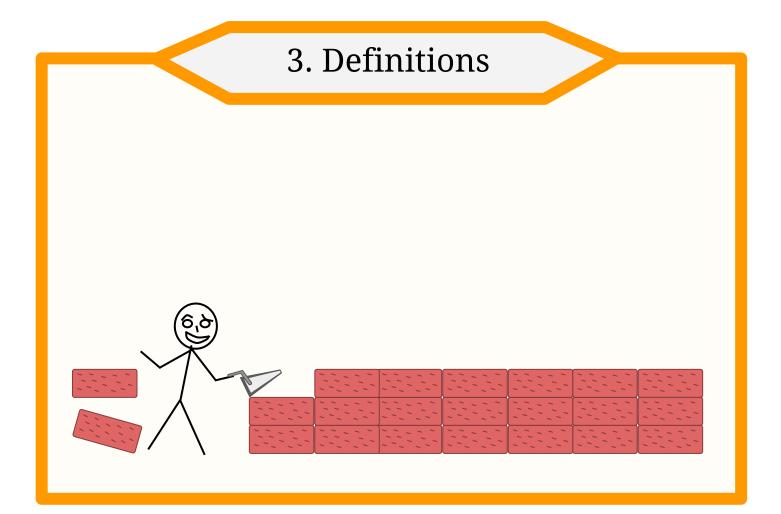


Savage: scoring rules $\leftarrow \rightarrow$ convex functions!



DECISION PROBLEMS $\leftarrow \rightarrow$ convex functions!





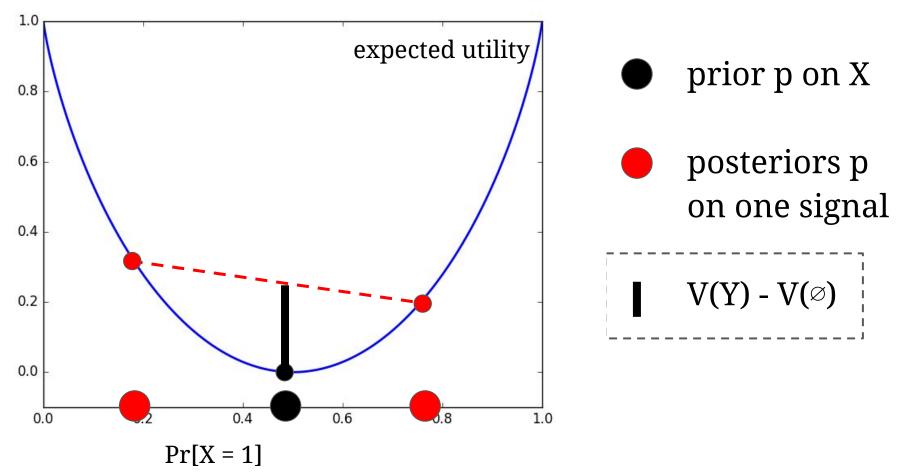
 $\label{eq:substitutes} \begin{array}{l} Y_1...Y_n \text{ are substitutes for } u \text{ if } V \text{ is submodular:} \\ \text{For } A \subseteq B \subseteq \{Y_1...Y_n\}, \end{array}$

$$V(A \cup \{Y_i\}) - V(A) \ge V(B \cup \{Y_i\}) - V(B).$$

- complements = supermodular
- depends on **both decision prob AND info structure**

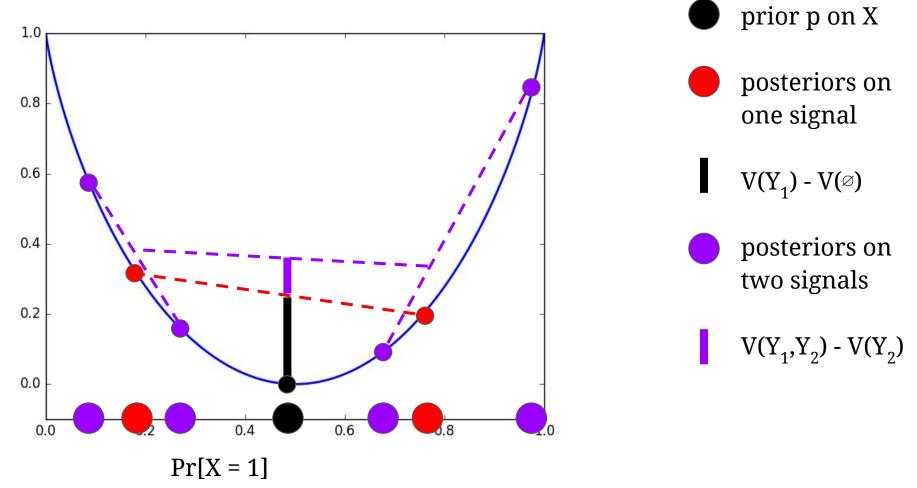
Visualizing an example of substitutes

Example: $S(q, x) = \log q(x)$. Y_1, Y_2 i.i.d. conditioned on x.

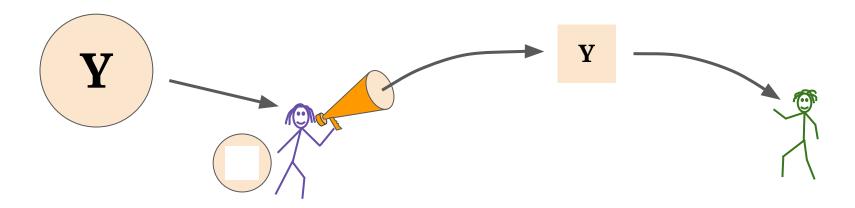


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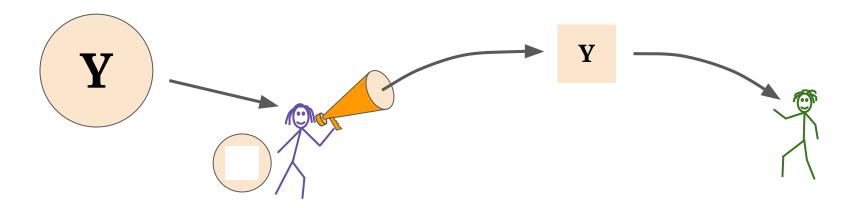


Roadblock: Information is divisible!



"Half the truth is often a great lie." - Benjamin Franklin

Roadblock: Information is divisible!



"Half the truth is often a great lie." - Benjamin Franklin

Example: Alice observes entire stock market, but strategically reports one stock's performance. $Y_1...Y_n$ are strong substitutes for u if: For $A \subseteq B \subseteq \{Y_1...Y_n\}$ and any randomized function f,

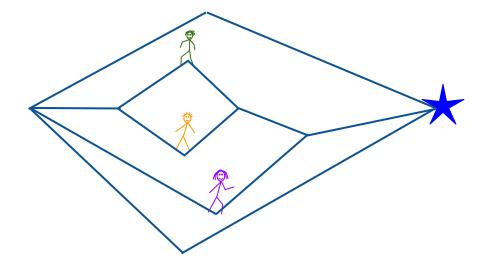
$$V(A \lor f(Y_i)) - V(A) \ge V(B \lor f(Y_i)) - V(B).$$

where A \vee f(Y_i) is the signal conveying both A and f(Y_i).

Lattice: partially ordered set closed under:

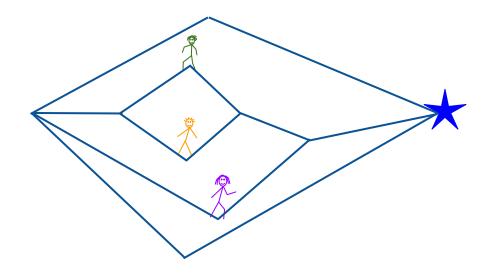
- meet $A \land B$: greatest C that is less than both
- **join** $A \lor B$: least C that is greater than both

Example: subsets. (meet = intersection, join = union) For signals: meet = "observe both", join = "common knowledge".



The "continuous" signal lattice

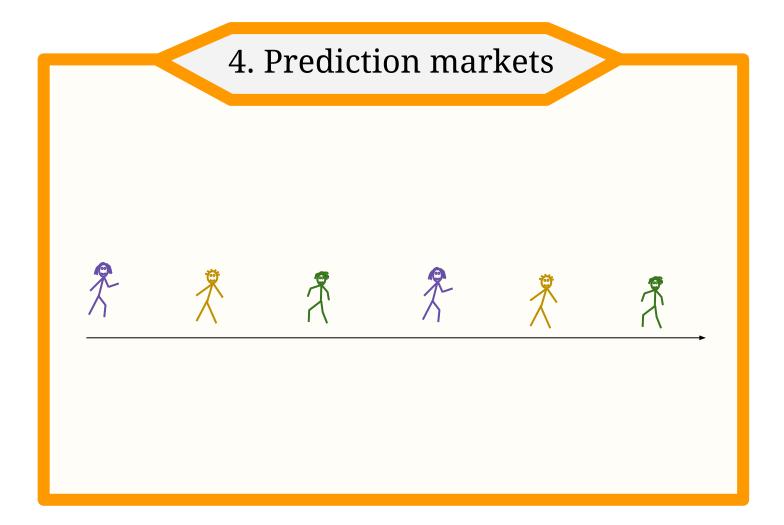
Idea: given X and Y₁ ... Y_n, The set of **randomized functions of sets of Y_is** form a lattice of signals ordered by informativeness (about X).



-- Remainders --

Two separate applications:

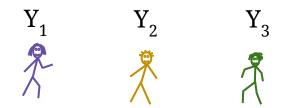
- Markets (for information). substitutes $\leftarrow \rightarrow$ good equilibria
- Algorithms. complexity of optimal info. acquisition



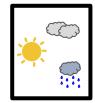
Idea / motivation

Each agent has a signal Y_i.

Goal: aggregate into prediction about X quickly.



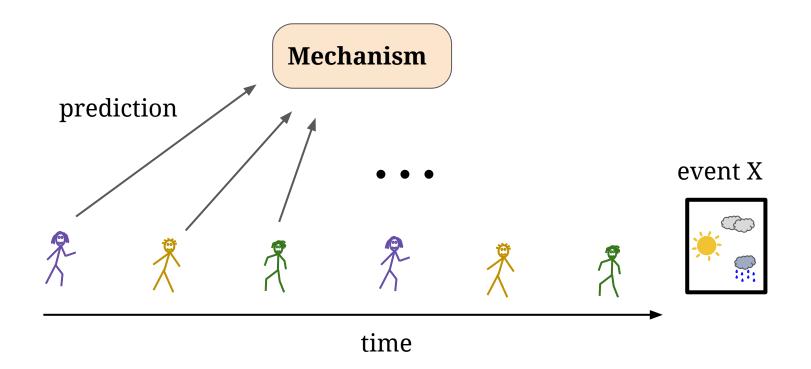
event X



Idea / motivation

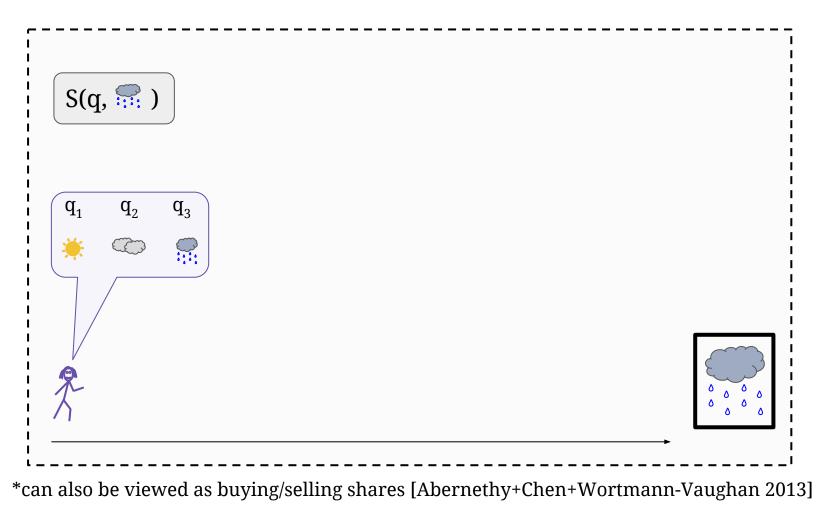
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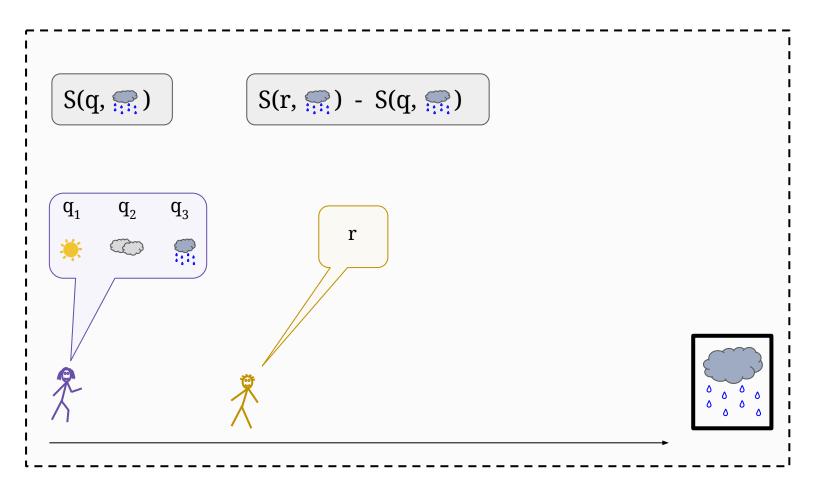
The mechanism [Hanson 2003]*

Only one participant: proper scoring rule! Truthful.



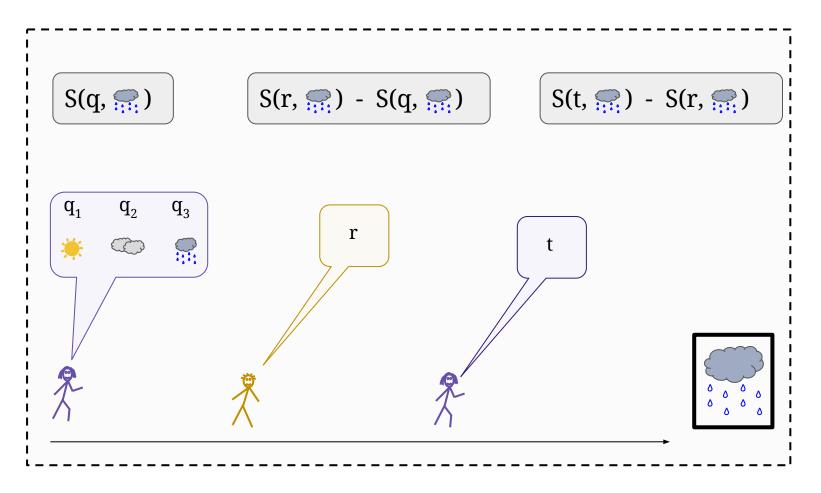
The mechanism [Hanson 2003]

Two participants: "chained" scoring rule! Truthful.



The mechanism [Hanson 2003]

Two participants, three stages: not understood!



The mechanism [Hanson 2003]

Two participants, three stages: **not understood!**

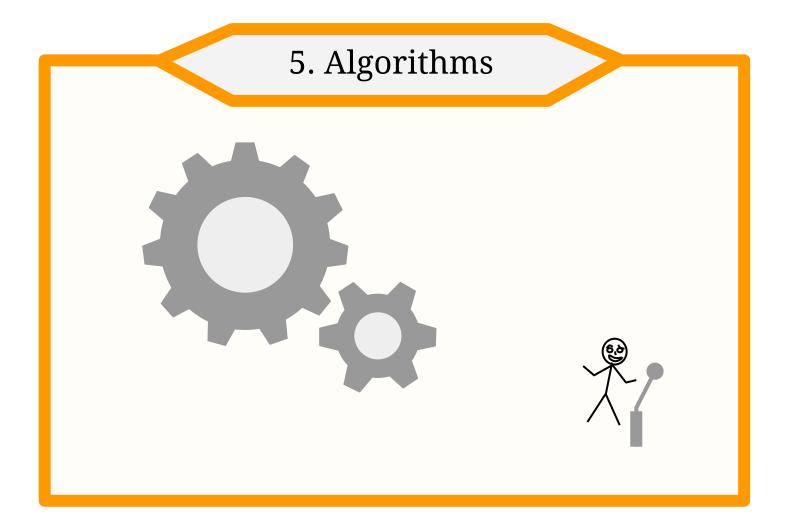


Thm. If and only if signals are strong **substitutes**, the only equilibria are "**all rush**".

(efficient market hypothesis $\leftarrow \rightarrow$ substitutes)

Thm. If and only if signals are strong **complements**, the only equilibria are "**all delay**".

(market failure $\leftarrow \rightarrow$ complements)



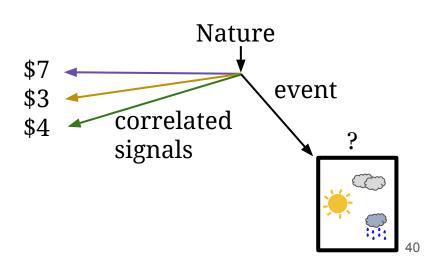
Algorithmic question "SignalSelection"

Input:

- utility function u (as an oracle...)
- joint distribution X, $Y_1 \dots Y_n$ (as an oracle...)
- prices $\pi_1 \dots \pi_n$ for the signals, budget constraint B

Output:

• which signals to acquire



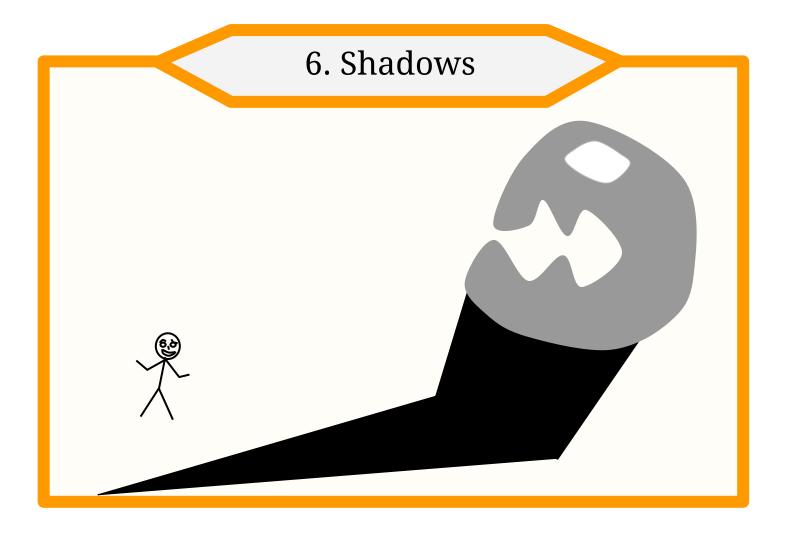
Complexity results

Reduction: SignalSelection \rightarrow set function maximization. Substitutes \Rightarrow 1-1/e approx in polynomial time

Reduction: set function maximization \rightarrow SignalSelection. Comps/generally \Rightarrow no approx w/ subexp. queries

Notes:

- As in submod. maximization, can handle e.g. matroid constraints.
- Ideas not new here at all! See survey [Krause+Guestrin 2011]
- Model / generality, focus of our question are new



What do we *know* about subs and comps?

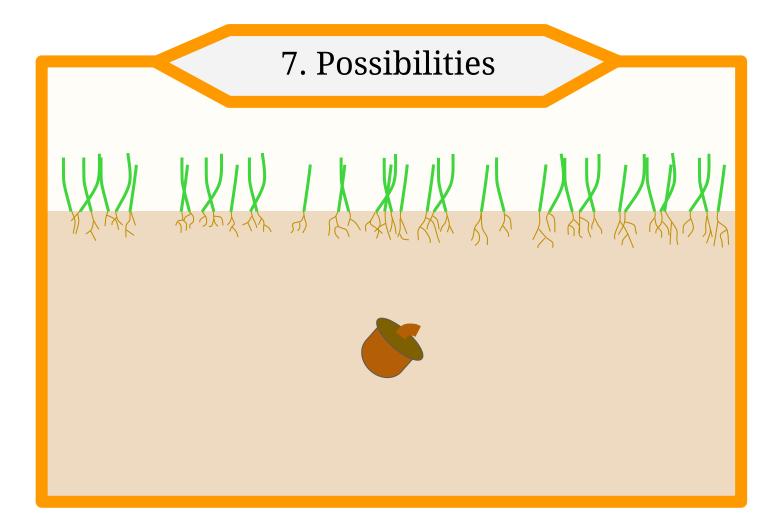
Examples: log scoring rule (Shannon entropy).

Intuition / geometry:

- high curvature in convex G $\, \leftarrow \, \rightarrow$ high marginal value of information
- to increase substitutability, put curvature close to the prior

"Universal" substitutes? - only with a "trivial-ish" structure

"Universal" complements? - Yes! XOR of bits



Open problems (small selection)

Game theory

- selling information
- signalling
- bundling complements
- other useful applications?

Algorithms

- check if signals are strong substitutes
- compute Alice's best response in stage one (decompose signal into sub. and comp. components)
- SignalSelection on discrete or continuous lattices!

Structure

- examples of (classes of) subs and comps
- "more substitutable" signal structures? utilities?
- "universal" substitutes and complements
- connections: e.g. sensitivity of Boolean functions?

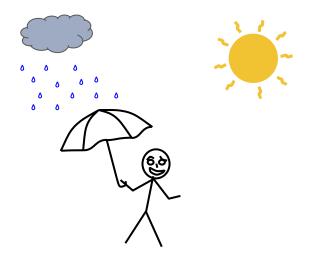
These slides:

bowaggoner.com/

Blog posts on proper scoring rules, generalized entropies, ...

bowaggoner.com/blog/

Information elicitation tutorial: sites.google.com/site/informationelicitation/



Thanks!